

Krasnoselskii-Mann fixed point iterations: sharp convergence rates

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Outline

- 1 Motivation & main results
- 2 Sharp recursive bounds by optimal transport
- 3 Random walks and Baillon-Bruck's conjecture
- 4 Optimal transports and the best possible constant
- 5 Discussion

Warmup: Banach-Picard iteration

When $T : X \rightarrow X$ is a ρ -contraction

(BP)

$$x^{n+1} = Tx^n$$

$$\|x^{n+1} - x^n\| = \|Tx^n - x^n\| \leq \rho^n \|Tx^0 - x^0\| \rightarrow 0$$

↓

convergence + error estimates + stopping rule

Krasnoselskii-Mann sequential averaging process

$T : C \rightarrow C$ non-expansive / C convex bounded in $(X, \|\cdot\|)$

(KM)

$$x^{n+1} = (1 - \alpha_{n+1}) x^n + \alpha_{n+1} T x^n$$

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- algorithm to compute & prove existence of fixed points
- **convex optimization:** Prox, Gradient, Douglas-Rachford, ADMM, POCS
- **dynamic programming:** stochastic shortest paths, Q-learning
- **continuous semigroups:** discretization of $\frac{dx}{dt} + [I - T](x) = 0$

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Questions:

- | | | |
|---|--------------------------------------|-------------------------|
| { | a) Existence of fixed points ? | |
| | b) Convergence of iterates ? | |
| | c) $\ T x^n - x^n\ \rightarrow 0$? | (Asymptotic Regularity) |
| | d) How fast ? | (Rate of Convergence) |

If we have $\|Tx^n - x^n\| \rightarrow 0$ then...

⇒ all strong/weak cluster points are fixed points of T

⇒ existence of fixed points: Browder-Göhde-Kirk'65

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and since $\|x^n - \bar{x}\|$ decreases for all $\bar{x} \in \text{Fix}T$

- ⇒ x^n converges strong/weak to a fixed point
- ⇒ convergence results: Krasnoselskii'55, Shaefer'57, Browder-Petryshyn'66, Edelstein'70, Groetsch'72, Ishikawa'76, Edelstein-O'Brien'78, Reich'79...

Example: θ -rotation in \mathbb{R}^2

$$T(x, y) = \begin{bmatrix} \cos \theta & -\sin \theta \\ \sin \theta & \cos \theta \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} \Rightarrow \|Tx^n - x^n\| = 2 \sin\left(\frac{\theta}{2}\right) \cos^n\left(\frac{\theta}{2}\right) \|x^0\|$$

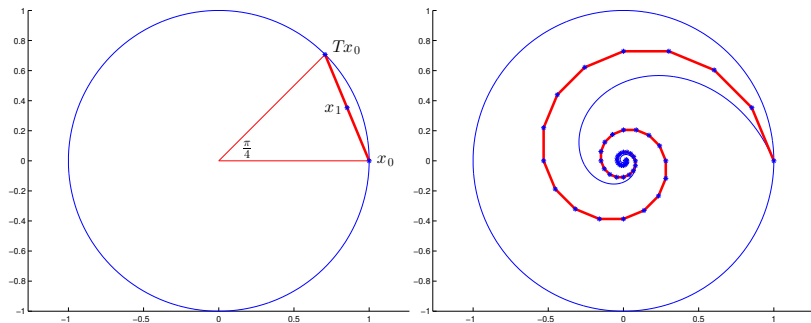


Figure: KM iterates for $\theta = \frac{\pi}{4}$ with $\alpha_n \equiv \frac{1}{2}$

Example: Right-shift on $\ell^1(\mathbb{N})$

$C = \{(p_i)_{i \in \mathbb{N}} \geq 0 : \sum_{i=0}^{\infty} p_i = 1\}$ is bounded with $\text{diam}(C) = 2$

$T(p_0, p_1, p_2, \dots) = (0, p_0, p_1, p_2, \dots)$ is an isometry

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$$p^0 = (1, 0, 0, 0, \dots)$$

$$p^1 = (1 - \alpha_1, \alpha_1, 0, 0, \dots)$$

$$p^2 = ((1 - \alpha_2)(1 - \alpha_1), (1 - \alpha_2)\alpha_1 + \alpha_2(1 - \alpha_1), \alpha_2\alpha_1, 0, \dots)$$

$$p^3 = \dots$$

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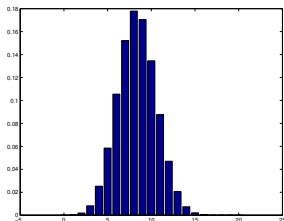
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$$p_k^n = \mathbb{P}(X_1 + \dots + X_n = k)$$

$$X_i \sim \text{Bernoulli}(\alpha_i)$$

$$\|Tp^n - p^n\|_1 = 2 \max_k p_k^n$$

60 years of $\|Tx^n - x^n\| \rightarrow 0$

- 1955 **Krasnoselskii**: X uniformly convex, C compact, $\alpha_n \equiv 1/2$
- 1957 **Schaefer**: extension to $\alpha_n \equiv \alpha$
- 1966 **Ishikawa**: extension to X strictly convex
- 1966 **Browder-Petryshyn**: X uniformly convex, $\text{Fix}(T) \neq \emptyset$, $\alpha_n \equiv \alpha$
- 1976 **Ishikawa**: X general, C compact, $\sum \alpha_n = \infty$, $\alpha_n \leq 1 - \epsilon$
- 1978 **Edelstein-O'Brien**: $\|Tx^n - x^n\| \rightarrow 0$ uniformly w.r.t x^0
- 1983 **Goebel-Kirk**: uniformly w.r.t. x^0 and T (for C given)
- 1992 **Baillon-Bruck**: $\alpha_n \equiv \alpha \Rightarrow \|Tx^n - x^n\| = O(1/\log n)$
- 1996 **Baillon-Bruck**: $\alpha_n \equiv \alpha \Rightarrow \|Tx^n - x^n\| = O(1/\sqrt{n})$
- 2003 **Kohlenbach**: $\|Tx^n - x^n\| \rightarrow 0$ only depends on $\text{diam}(C)$

Conjecture (Baillon-Bruck 1992)

There exists a universal constant κ such that

$$\|x^n - Tx^n\| \leq \kappa \frac{\text{diam}(C)}{\sqrt{\sum_{i=1}^n \alpha_i(1-\alpha_i)}}. \quad (\text{BB})$$

Theorem (Baillon-Bruck 1996)

For $\alpha_n \equiv \alpha$ constant, (BB) holds with $\kappa = 1/\sqrt{\pi}$.

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For $\alpha_n \equiv \alpha$ constant, (BB) holds with $\kappa = 1/\sqrt{\pi}$.

Our contributions (C-Soto-Vaisman 2014, Bravo-C 2016, Bravo-C-Pavez 2017)

- 1 (BB) holds for general α_n with $\kappa = 1/\sqrt{\pi} \sim 0.5642$
- 2 Nonlinear maps: the constant $\kappa = 1/\sqrt{\pi}$ is the best possible
- 3 Affine maps: tight bound with $\kappa = \max_z \sqrt{z}e^{-z}I_0(z) \sim 0.4688$
- 4 Inexact KM: $x^{n+1} = (1-\alpha_{n+1})x^n + \alpha_{n+1}(Tx^n + e^{n+1})$

$$\|x^n - Tx^n\| \leq \kappa \frac{\text{diam}(C)}{\sqrt{\sum_{i=1}^n \alpha_i(1-\alpha_i)}} \quad (\text{BB})$$

In this talk...

- 1 (BB) holds with $\kappa = 1/\sqrt{\pi} \sim 0.5642$
- 2 The constant $\kappa = 1/\sqrt{\pi}$ is the best possible

Along the way we unveil the probabilistic structure behind (KM)...

and discover that Krasnoselskii was “almost” right in taking $\alpha_n \equiv \frac{1}{2}$.

Streamline

Algorithms for convex optimization



Fixed point iterations



Recursive bounds by optimal transport



Markov processes and combinatorics



Rates of convergence with optimal constants



Selection of optimal parameters

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A recursive bound $\|x^m - x^n\| \leq d_{mn}$

(KM)

$$x^{n+1} = (1 - \alpha_{n+1})x^n + \alpha_{n+1}Tx^n$$

- We look for sharp bounds for $\|Tx^n - x^n\| = \|x^{n+1} - x^n\| / \alpha_{n+1}$
- We achieve this by bounding $\|x^m - x^n\| \leq d_{mn}$ for all $m \leq n$

A recursive bound $\|x^m - x^n\| \leq d_{mn}$

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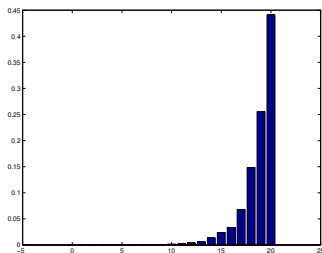
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Using the conventions $\alpha_0 = 1$ and $Tx^{-1} = x^0$, we have

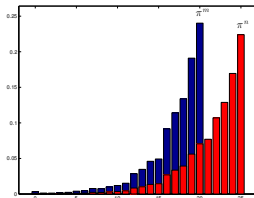
$$x^n = \sum_{i=0}^n \pi_i^n Tx^{i-1}$$

$$\pi_i^n = \alpha_i \prod_{k=i+1}^n (1 - \alpha_k)$$



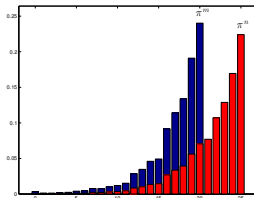
A recursive bound $\|x^m - x^n\| \leq d_{mn}$

$$x^m - x^n = \sum_{i=0}^m \pi_i^m T x^{i-1} - \sum_{j=0}^n \pi_j^n T x^{j-1}$$



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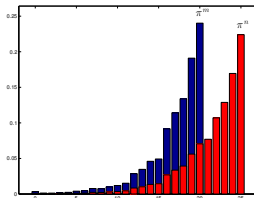
Let P_{mn} be the set of transport plans $z \geq 0$ taking π^m to π^n

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A recursive bound $\|x^m - x^n\| \leq d_{mn}$

$$x^m - x^n = \sum_{i=0}^m \sum_{j=0}^n z_{ij} T x^{i-1} - \sum_{j=0}^n \sum_{i=0}^m z_{ij} T x^{j-1}$$



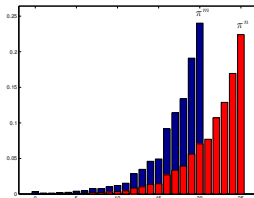
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$$x^m - x^n = \sum_{i=0}^m \sum_{j=0}^n z_{ij} [Tx^{i-1} - Tx^{j-1}]$$

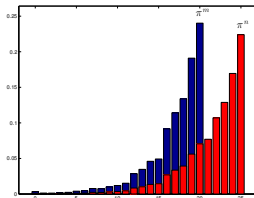


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A recursive bound $\|x^m - x^n\| \leq d_{mn}$

$$\|x^m - x^n\| \leq \sum_{i=0}^m \sum_{j=0}^n z_{ij} \|x^{i-1} - x^{j-1}\|$$

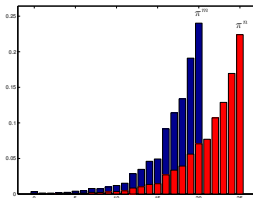


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$$\|x^m - x^n\| \leq \sum_{i=0}^m \sum_{j=0}^n z_{ij} d_{i-1, j-1}$$

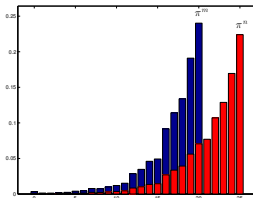


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$$\|x^m - x^n\| \leq \sum_{i=0}^m \sum_{j=0}^n z_{ij} d_{i-1,j-1} \quad \longrightarrow \quad \min_z$$



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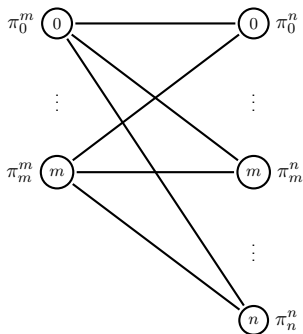
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Recursive optimal transports

Assume w.l.o.g. that $\text{diam}(C)=1$ and set $d_{-1,n} = 1$ for all $n \in \mathbb{N}$

Define d_{mn} inductively for $0 \leq m \leq n$ as

$$(R) \quad d_{mn} \triangleq \min_{z \in P_{mn}} \sum_{i=0}^m \sum_{j=0}^n z_{ij} d_{i-1,j-1}$$



Property

Let $T : C \rightarrow C$ be a nonexpansive map on a convex set C with $\text{diam}(C) = 1$. Then the KM iterates satisfy $\|x^m - x^n\| \leq d_{mn}$ for all $m, n \in \mathbb{N}$.

The bounds d_{mn} are fully determined by the α_k 's and are:

- Independent of the space X
- Independent of the operator $T : C \rightarrow C$

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Theorem (Bravo-C 2016)

These bounds are the best possible: There exists a nonexpansive $T : C \rightarrow C$ on the unit cube $C = [0, 1]^{\mathbb{N}} \subseteq \ell^\infty(\mathbb{N})$ and a corresponding KM sequence such that $\|x^m - x^n\|_\infty = d_{mn}$ for all $m, n \in \mathbb{N}$.

Proof: T built from dual solutions of the optimal transports. □

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⇒ The d_{mn} 's are the right object to obtain sharp bounds !

Restatement of (BB)

$$\|Tx^n - x^n\| = \left\| \frac{x^{n+1} - x^n}{\alpha_{n+1}} \right\| \leq \frac{d_{n,n+1}}{\alpha_{n+1}} = ?$$

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$$\frac{d_{n,n+1}}{\alpha_{n+1}} \leq \frac{1}{\sqrt{\pi}} \frac{1}{\sqrt{\sum_{k=1}^n \alpha_k (1 - \alpha_k)}} \quad ?$$

Metric properties of the optimal transport bounds

Theorem (Aygen-Satik 1994, Bravo-C 2016)

$d(m, n) = d_{mn}$ defines a metric over the set $\mathcal{N} = \{-1, 0, 1, 2, 3, \dots\}$.

Original proof is 30+ pages long. Short proof in Bravo-C 2016 (3 pages).

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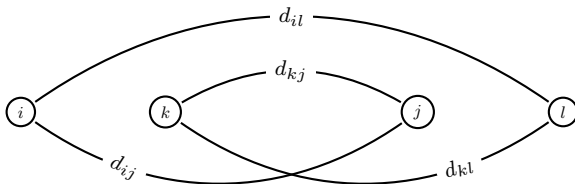
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Theorem (Aygen-Satik 1994, Bravo-C. 2016)

For all $i \leq k \leq j \leq l$ we have $d_{il} + d_{kj} \leq d_{ij} + d_{kl}$.

Original proof is 80+ pages long but holds for $\alpha_n \in [0, 1]$.

We found a 2-page proof for the case $\alpha_n \geq 1/2$.



Optimal transport for $\alpha_n \geq 1/2$

$$\pi_0^m \textcircled{0}$$

$$\textcircled{0} \pi_0^n$$

$$\pi_1^m \textcircled{1}$$

$$\textcircled{1} \pi_1^n$$

$$\pi_2^m \textcircled{2}$$

$$\textcircled{2} \pi_2^n$$

$$\vdots$$

$$\vdots$$

$$\pi_m^m \textcircled{m}$$

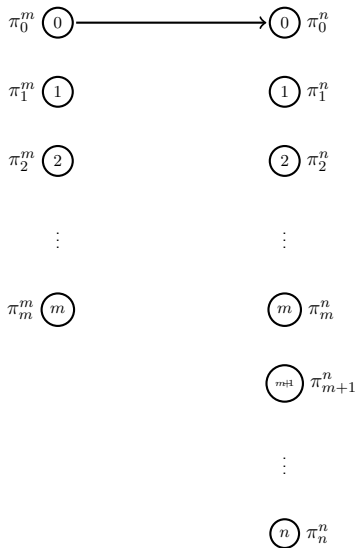
$$\textcircled{m} \pi_m^n$$

$$\textcircled{m+1} \pi_{m+1}^n$$

$$\vdots$$

$$\textcircled{n} \pi_n^n$$

Optimal transport for $\alpha_n \geq 1/2$



Optimal transport for $\alpha_n \geq 1/2$

$$\pi_0^m \textcircled{0} \longrightarrow \textcircled{0} \pi_0^n$$

$$\pi_1^m \textcircled{1} \longrightarrow \textcircled{1} \pi_1^n$$

$$\pi_2^m \textcircled{2} \qquad \qquad \textcircled{2} \pi_2^n$$

$$\vdots$$

$$\vdots$$

$$\pi_m^m \textcircled{m}$$

$$\textcircled{m} \pi_m^n$$

$$\textcircled{m+1} \pi_{m+1}^n$$

$$\vdots$$

$$\textcircled{n} \pi_n^n$$

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$$\vdots$$

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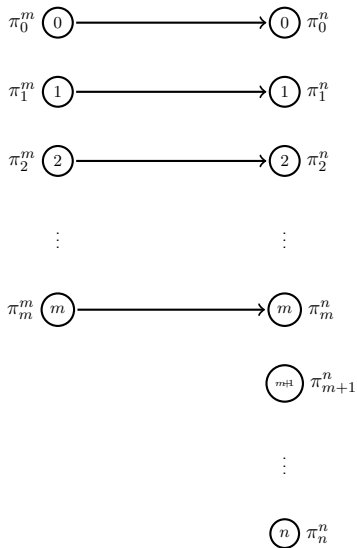
$$\textcircled{m} \pi_m^n$$

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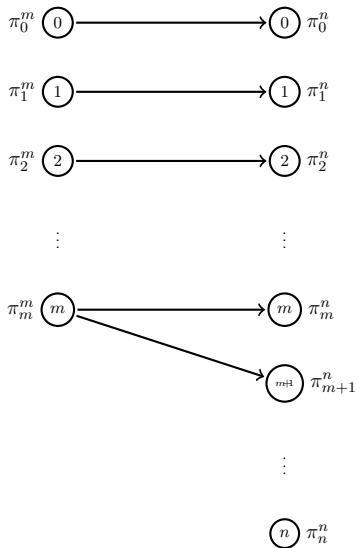
$$\vdots$$

$$\textcircled{n} \pi_n^n$$

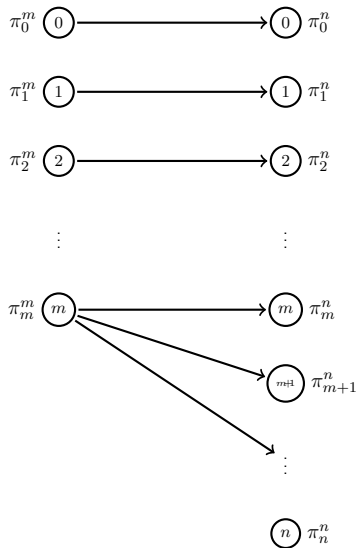
Optimal transport for $\alpha_n \geq 1/2$



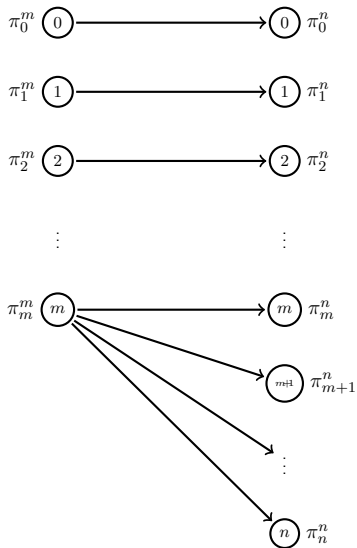
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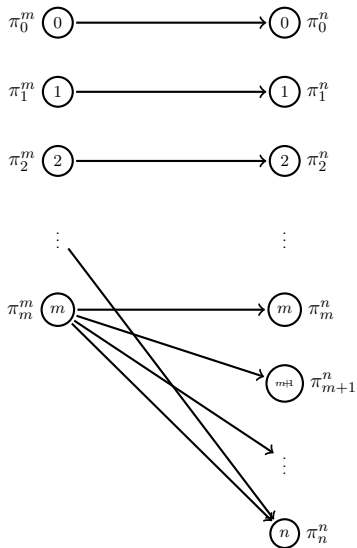
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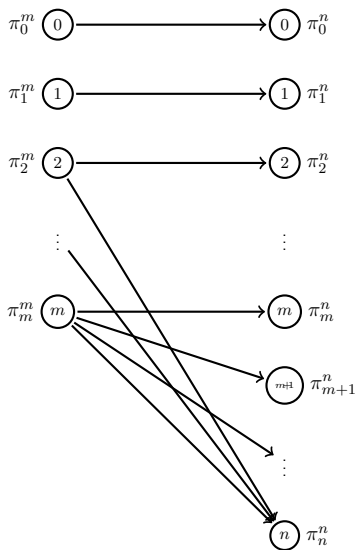
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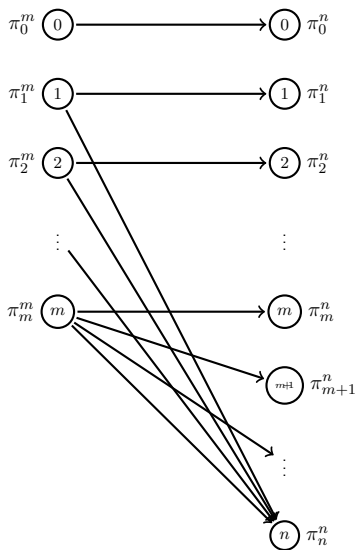
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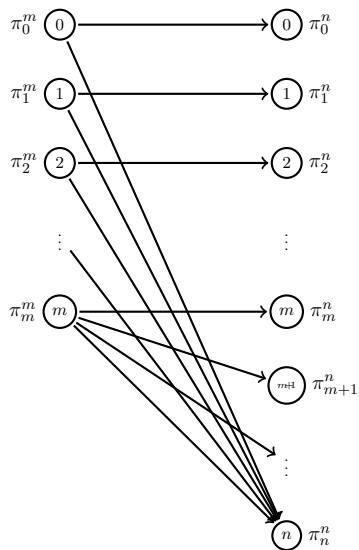
Optimal transport for $\alpha_n \geq 1/2$



Optimal transport for $\alpha_n \geq 1/2$



Optimal transport for $\alpha_n \geq 1/2$



Optimal transport for $\alpha_n \geq 1/2$

The optimal solution is

$$\begin{aligned}
 z_{ii} &= \pi_i^n && \text{for } i = 0, \dots, m, \\
 z_{mj} &= \pi_j^n && \text{for } j = m+1, \dots, n-1, \\
 z_{in} &= \pi_i^m - \pi_i^n && \text{for } i = 0, \dots, m-1, \\
 z_{mn} &= \pi_m^m - \sum_{j=m}^{n-1} \pi_j^n.
 \end{aligned}$$

and the recursion “*simplifies*” to

$$d_{mn} = \sum_{j=m+1}^{n-1} \pi_j^n d_{m-1, j-1} + \sum_{i=0}^{m-1} (\pi_i^m - \pi_i^n) d_{i-1, n-1} + (\pi_m^m - \sum_{j=m}^{n-1} \pi_j^n) d_{m-1, n-1}$$

$$\alpha_n \equiv \alpha \geq \frac{1}{2}$$

$$\begin{aligned}
 d_{6,10}(\alpha) = & \alpha(4 - 36\alpha + 328\alpha^2 - 2671\alpha^3 + 19853\alpha^4 - 132880\alpha^5 + 785003\alpha^6 \\
 & - 4016624\alpha^7 + 17541102\alpha^8 - 64796454\alpha^9 + 201809157\alpha^{10} \\
 & - 530670200\alpha^{11} + 1183318617\alpha^{12} - 2250818306\alpha^{13} + 3675506816\alpha^{14} \\
 & - 5184593492\alpha^{15} + 6352439437\alpha^{16} - 6792441644\alpha^{17} + 6361687020\alpha^{18} \\
 & - 5232669869\alpha^{19} + 3785701567\alpha^{20} - 2409974375\alpha^{21} + 1348858198\alpha^{22} \\
 & - 662337623\alpha^{23} + 284299971\alpha^{24} - 106102624\alpha^{25} + 34171973\alpha^{26} \\
 & - 9400913\alpha^{27} + 2178730\alpha^{28} - 417352\alpha^{29} + 64328\alpha^{30} \\
 & - 7667\alpha^{31} + 663\alpha^{32} - 37\alpha^{33} + \alpha^{34})
 \end{aligned}$$

...??????????

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Upper estimate: $d_{mn} \leq c_{mn}$

Consider the feasible non-optimal transport plan

$$\hat{z}_{ij} = \begin{cases} \pi_j^n & \text{for } i = j \leq m \\ \pi_i^m \pi_j^n & \text{for } j = m + 1, \dots, n \\ 0 & \text{otherwise} \end{cases}$$

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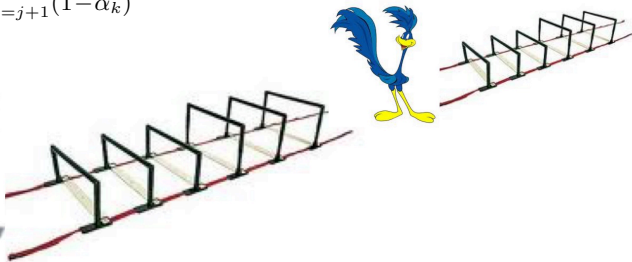
Setting $c_{-1,n} = 1$ we get inductively

$$\|x^m - x^n\| \leq d_{mn} \leq c_{mn} \triangleq \sum_{i=0}^m \sum_{j=m+1}^n \pi_i^m \pi_j^n c_{i-1,j-1}$$

Probabilistic interpretation of c_{mn}

$$\mathbb{P}[C_i = 1] = \mathbb{P}[R_i = 1] = \alpha_i$$

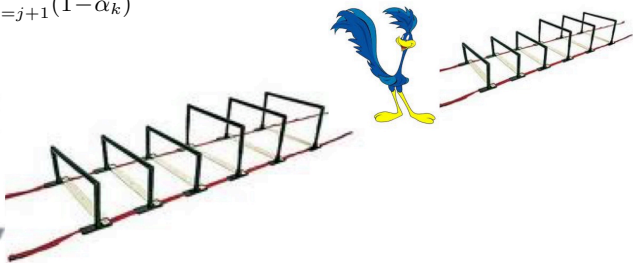
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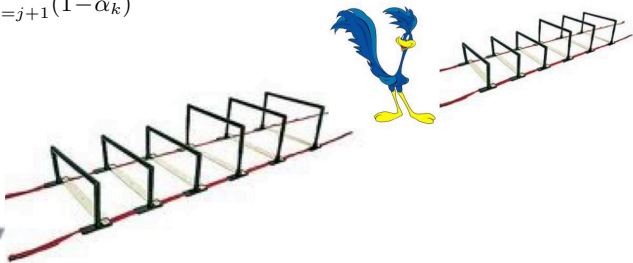


$$\tilde{c}_{mn} = \mathbb{P}[\text{roadrunner escapes}]$$

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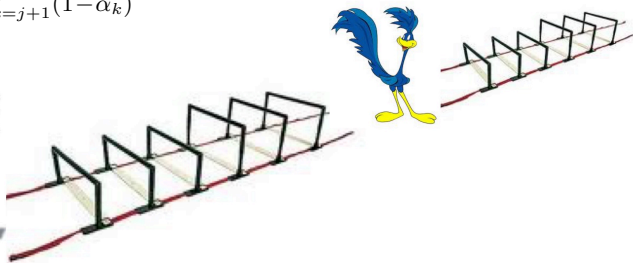


$$\tilde{c}_{mn} = \sum_{i=0}^m \sum_{j=m+1}^n \pi_i^m \pi_j^n \tilde{c}_{i-1, j-1}$$

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$$c_{mn} = \mathbb{P} \left(\sum_k^n C_i > \sum_k^m R_i, \forall k = m + 1, \dots, 1 \right)$$

Coyote must fall more often than Roadrunner

The random walk and the gambler's ruin appear...

Using probabilistic arguments we get an explicit formula for the bound

$$\|Tx^n - x^n\| \leq \frac{c_{n,n+1}}{\alpha_{n+1}} = \mathbb{E}[F(M)]$$

where M is a sum of independent nonhomogeneous Bernoulli trials

$$M = M_1 + \dots + M_n \quad ; \quad \mathbb{P}(M_i=1) = p_i \triangleq 2\alpha_i(1-\alpha_i)$$

and $F(m)$ is the probability that a standard random walk in \mathbb{Z} remains non-negative for the first m stages

$$F(m) = \binom{m}{\lfloor m/2 \rfloor} 2^{-m}.$$

Upper estimate

Thus (BB) has been reduced to show that

$$\mathbb{E}[F(M)] \leq \frac{1}{\sqrt{\pi \sum_{i=1}^n \alpha_i (1 - \alpha_i)}}$$

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$$\underbrace{\sqrt{\frac{\pi}{2} (p_1 + \dots + p_n)} \mathbb{E}[F(M_1 + \dots + M_n)]}_{R_n(p)} \leq 1$$

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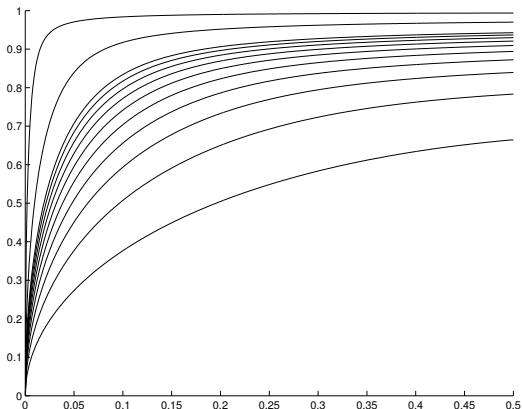
$$\underbrace{\sqrt{\frac{\pi}{2}(p_1 + \dots + p_n)} \mathbb{E}[F(M_1 + \dots + M_n)]}_{R_n(p)} \leq 1$$

Lemma

$R_n(p)$ is maximal when $p_i \in \{u, \frac{1}{2}\}$ for some $0 < u < \frac{1}{2}$

Upper estimate – Case 1: all $p_i = u$

$$R_n(p) = \sqrt{\frac{\pi}{2}nu} \mathbb{E}[F(B(n, u))] = \sqrt{\frac{\pi}{2}nu} {}_2F_1\left(-n, \frac{1}{2}; 2; 2u\right) \leq 1$$



Upper estimate – Case 2: some $p_i = \frac{1}{2}$

Suppose $p_1 = \frac{1}{2}$ and let $S = M_2 + \dots + M_n$. Conditioning on M_1

$$\mathbb{E}[F(M)] = \mathbb{E}[G(S)]$$

where $G(k) = \frac{1}{2}[F(k) + F(k+1)]$ is convex so we may use the following Hoeffding-type inequality

Theorem (C.-Soto-Vaisman'2014, *Israel J. Math.*)

$$\mathbb{E}[G(S)] \leq \mathbb{E}[G(Z)] \text{ where } Z \sim \text{Poisson}(z) \text{ with } z = \mathbb{E}(S).$$

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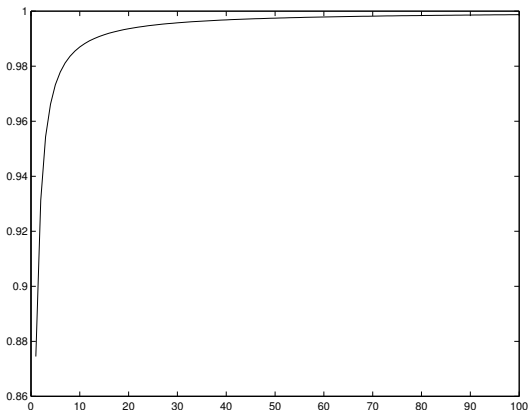
$$\mathbb{E}[G(S)] \leq \mathbb{E}[G(Z)] \text{ where } Z \sim \text{Poisson}(z) \text{ with } z = \mathbb{E}(S).$$

$$\Rightarrow \mathbb{E}[F(M)] \leq \mathbb{E}[G(Z)] = I_0(z) + (1 - \frac{1}{2z})I_1(z)$$

with $I_0(z), I_1(z)$ modified Bessel functions

Upper estimate – Case 2: some $p_i = \frac{1}{2}$

$$R(p) \leq \sqrt{\frac{\pi}{2}(\frac{1}{2} + z)} [I_0(z) + (1 - \frac{1}{2z})I_1(z)] \leq 1$$



Conclusion: (BB) holds with $\kappa = 1/\sqrt{\pi} \sim 0.5642$

Both cases combined yield

Theorem (C.-Soto-Vaisman'2014, *Israel J. Math.*)

$$\|Tx^n - x^n\| \leq \frac{\text{diam}(C)}{\sqrt{\pi \sum_{k=1}^n \alpha_k (1 - \alpha_k)}}$$

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Remarks: (Baillon-C.-Vaisman'2016, *Comb. Prob. & Computing*)

- If T is affine (BB) holds with $\kappa = \max_{u \geq 0} \sqrt{u} e^{-u} I_0(u) \sim 0.4688$.
- This bound is tight and is attained by the right shift of Example 2.

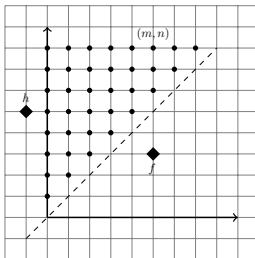
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A related stochastic process

$$d_{mn} = \sum_{i=0}^m \sum_{j=0}^n z_{ij}^{mn} d_{i-1,j-1}$$

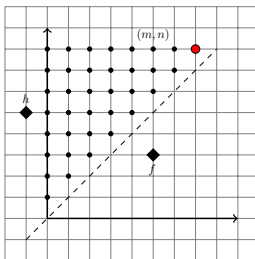
d_{mn} can be interpreted as the absorption probability of a finite Markov chain \mathcal{D} with state space $\mathcal{S} = \{(m, n) : 0 \leq m < n\} \cup \{h, f\}$.



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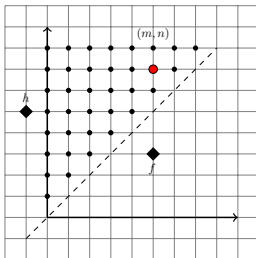
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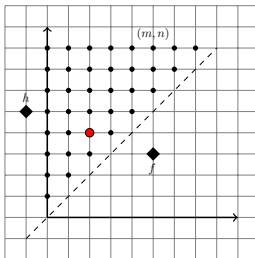
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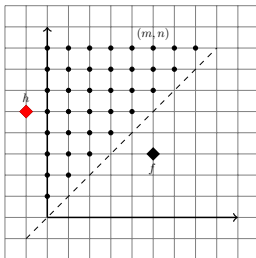
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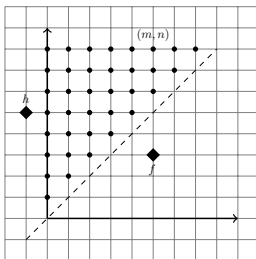
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Similarly, for the Markov chain \mathcal{C} defined by the sub-optimal transports

$$c_{mn} = \sum_{i=0}^m \sum_{j=m+1}^n \pi_i^m \pi_j^n c_{i-1,j-1} = \mathbb{P}(\mathcal{C} \text{ attains } h|mn)$$

Baillon-Bruck bound for constant $\alpha_n \equiv \alpha \in [1/2, 1[$

$$\sqrt{n\alpha(1-\alpha)} \|x^n - Tx^n\| \leq \kappa ?$$

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where

$$\kappa_n(\alpha) = \sqrt{n\alpha(1-\alpha)} d_{n,n+1}(\alpha)/\alpha,$$

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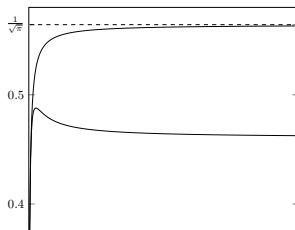
$$\kappa_n(\alpha) = \sqrt{n\alpha(1-\alpha)} d_{n,n+1}(\alpha)/\alpha,$$

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Moreover, the previous analysis of the c_{mn} 's gives explicitly

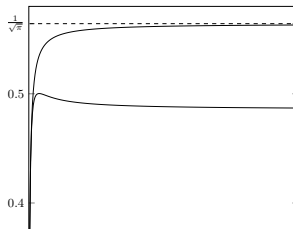
$$\tilde{\kappa}_n(\alpha) = \frac{1}{\pi} \int_0^{4n\alpha(1-\alpha)} \sqrt{\frac{1}{s} - \frac{1}{4n\alpha(1-\alpha)}} \left(1 - \frac{s}{n}\right)^n ds \nearrow \frac{1}{\pi} \Gamma\left(\frac{1}{2}\right) = \frac{1}{\sqrt{\pi}}.$$

Difference between $\kappa_n(\alpha)$ and $\tilde{\kappa}_n(\alpha)$ for $n = 1, \dots, 300$



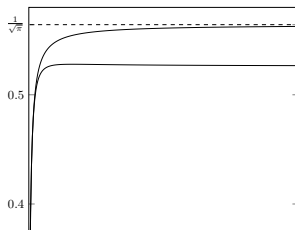
$\alpha = 0.5.$

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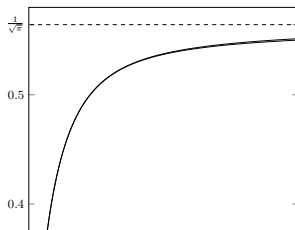
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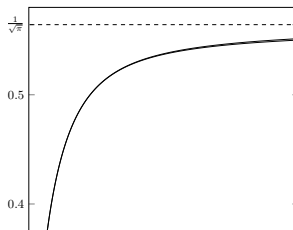
$\alpha = 0.85$.

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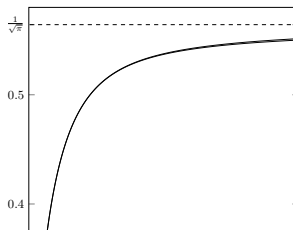


$\alpha = 0.99.$

Lemma

For $\alpha_n \equiv \alpha \in [\frac{1}{2}, 1[$ we have $0 \leq c_{n,n+1} - d_{n,n+1} \leq 4n(1-\alpha)^2$.

Difference between $\kappa_n(\alpha)$ and $\tilde{\kappa}_n(\alpha)$ for $n = 1, \dots, 300$



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Theorem (Bravo-C.'2016)

The constant $\kappa = \frac{1}{\sqrt{\pi}}$ is the best possible.

Proof. Taking $\alpha_n = 1 - \ln(n)/n$ we get $\kappa_n(\alpha_n) \sim \tilde{\kappa}_n(\alpha_n) \rightarrow \frac{1}{\sqrt{\pi}}$. □

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Best constant stepsize $\alpha_n \equiv \alpha$?

$$\|Tx^n - x^n\| \leq \gamma(\alpha)/\sqrt{n}$$

$$\gamma(\alpha) = \sup_{n \in \mathbb{N}} \sqrt{n} d_{n,n+1}(\alpha)/\alpha$$

The minimum seems to be attained at $\alpha = 0.4623\dots!$

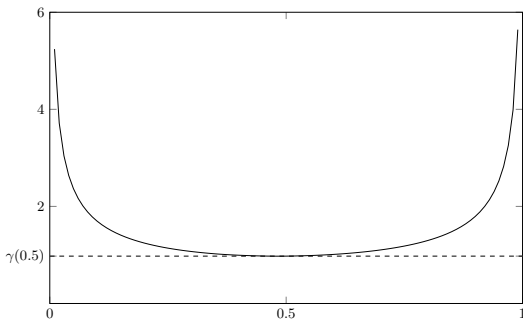


Figure: The rate $\gamma(\alpha)$ as a function of α (for $0.01 \leq \alpha \leq 0.99$).

Krasnoselskii's original iteration $\alpha_n \equiv \frac{1}{2}$

$$\gamma(\alpha) = \sup_{n \in \mathbb{N}} \sqrt{n} d_{n,n+1}(\alpha) / \alpha$$

For $\alpha = \frac{1}{2}$ the sup [seems](#) to be attained at $n = 8$

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$$\begin{aligned} d_{8,9}(\alpha) / \alpha &= 1 - 8\alpha + 64\alpha^2 - 448\alpha^3 + 2835\alpha^4 - 16008\alpha^5 + 79034\alpha^6 \\ &\quad - 334908\alpha^7 + 1201873\alpha^8 - 3622324\alpha^9 + 9129380\alpha^{10} \\ &\quad - 19214722\alpha^{11} + 33796129\alpha^{12} - 49776610\alpha^{13} + 61566687\alpha^{14} \\ &\quad - 64152608\alpha^{15} + 56488500\alpha^{16} - 42133404\alpha^{17} + 26651679\alpha^{18} \\ &\quad - 14288252\alpha^{19} + 6472429\alpha^{20} - 2462126\alpha^{21} + 778478\alpha^{22} \\ &\quad - 201354\alpha^{23} + 41584\alpha^{24} - 6604\alpha^{25} + 758\alpha^{26} - 56\alpha^{27} + 2\alpha^{28} \end{aligned}$$

\Rightarrow the sharp rate in Krasnoselskii's iteration [would be](#)

$$\gamma\left(\frac{1}{2}\right) = \frac{46302245}{67108864} \sqrt{2} \sim 0.9757 \quad \left(\text{smaller than } \frac{2}{\sqrt{\pi}} \sim 1.1284\right)$$

Sharp rate in Hilbert spaces?

- For $\alpha_n \equiv \alpha$, Browder & Petryshin proved $\sum \|x^{n+1} - x^n\|^2 < \infty$
- Since $\|x^{n+1} - x^n\|$ is decreasing, this readily gives

$$\|x^n - Tx^n\| = o(1/\sqrt{n})$$

- Which is the exact rate of asymptotic regularity in the Hilbert case?

What about errors?

Inexact KM:

$$x^{n+1} = (1 - \alpha_{n+1})x^n + \alpha_{n+1}(Tx^n + e^{n+1})$$

- **Controlled errors:**

Recently solved with M. Bravo and M. Pavez-Signé (available on arXiv).

- **Stochastic errors:**



Thanks !

Preprints available at

<https://sites.google.com/site/cominettiroberto/>