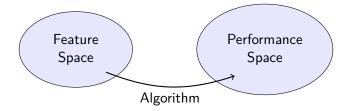
# Generating mixed integer programming instances with challenging properties

Simon Bowly

University of Melbourne

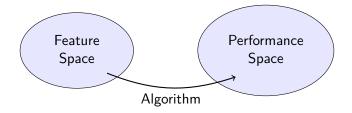
AMSI Optimise June 20, 2019

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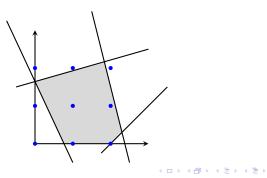
Give an overview of...

- MIP solvers (in particular branching strategies)
- Generating new instances using evolutionary algorithms
- Characteristics of the generated instances

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#### Mixed Integer Program

 $\begin{array}{ll} \text{maximize} & c^{\mathsf{T}}x\\ \text{subject to} & Ax \leq b\\ & x \geq 0\\ \text{some or all } x_i \text{ integral} \end{array}$ 



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Many interacting components make up a MIP solver:

- Presolvers
- Cutting planes
- Primal heuristics
- Parallelisation

- LP solvers
- Branching rules
- Node selection rules
- Domain propagation

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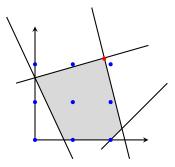
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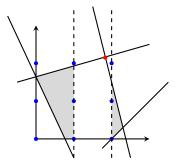
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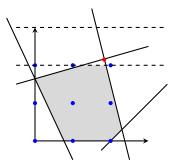
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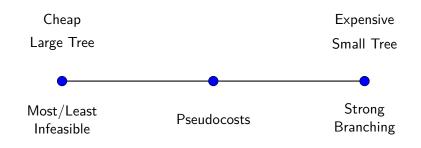
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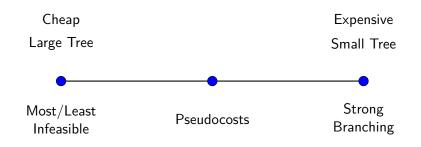
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To conduct a branching study  $\ldots$ 

- Disable cuts beyond the root node
- Provide the optimal solution

(Linderoth and Savelsbergh, 1999)

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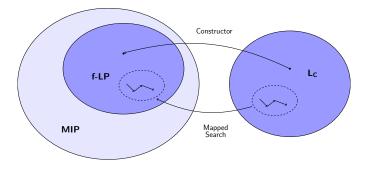
#### Genetic algorithms

- Population based metaheuristic with fitness selection
- Crossover is the key search operator

Previous applications to instance generation

- Finding worst-case bounds for sorting algorithm performance (Cotta and Moscato, 2003).
- Exposing strengths and weaknesses of heuristics (van Hemert, 2006; Langdon and Poli, 2007).

How can we narrow the search space?



Restrict search to feasible, bounded problems.

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#### **Dual programs**

$$\begin{array}{lll} \max & c^T x & \min & b^T y \\ \text{s.t.} & Ax + s = b & \Longleftrightarrow & \text{s.t.} & A^T y - r = c \\ & x, s \geq 0 & y, r \geq 0 \end{array}$$

#### **Optimality conditions**

$$(x, s)$$
 and  $(y, r)$  feasible to primal and dual  
 $x_i r_i = 0 \quad \forall i$   
 $y_j s_j = 0 \quad \forall j$ 

Construction

$$b = Ax + s$$
  
$$c = A^T y - r$$

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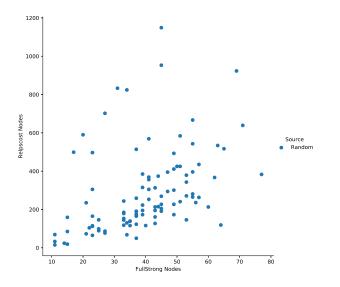
• Uniform row crossover:

$$\begin{bmatrix} a_{11} & \dots & a_{1n} \\ a_{21} & \dots & a_{2n} \\ a_{31} & \dots & a_{3n} \end{bmatrix} \begin{bmatrix} b_{11} & \dots & b_{1n} \\ b_{21} & \dots & b_{2n} \\ b_{31} & \dots & b_{3n} \end{bmatrix}$$
$$\begin{bmatrix} a_{21} & \dots & a_{2n} \\ b_{21} & \dots & b_{2n} \\ a_{31} & \dots & a_{3n} \end{bmatrix} \begin{bmatrix} a_{11} & \dots & a_{1n} \\ b_{31} & \dots & b_{3n} \\ b_{11} & \dots & b_{1n} \end{bmatrix}$$

- Multi-objective NSGA-II algorithm (Deb et al., 2002)
- Small instances 50 integer variables and 50 constraints
- fullstrong and relpscost branching strategies in SCIP 6.0.0

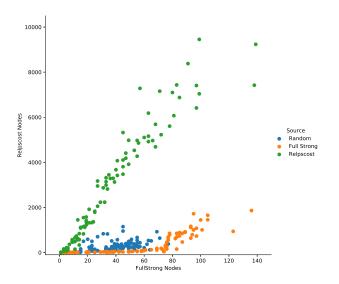
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## Performance Discriminating Instances



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## Performance Discriminating Instances

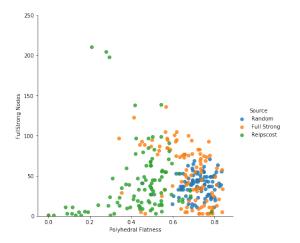


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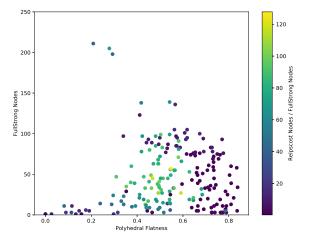
## Properties of Generated Instances



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## Properties of Generated Instances



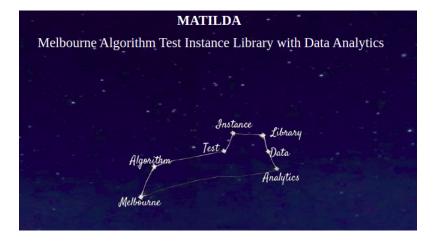
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- Evolutionary generation allows us to search performance space by manipulating instances.
- Small instances can be generated which expose performance differences in fundamental components.
- Branching is just one facet of MIP solver performance there's more of the performance space to explore.

## References

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## MATILDA



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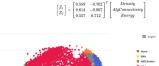
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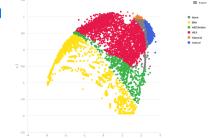
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#### Footprint Analysis

GRAPH COLORING PROBLEM

INSTANCE SPACE 2D COORDINATES = PROJECTION MATRIX X FEATURES VECTOR







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